International Journal of Scientific Engineering and Research (IJSER)
www.ijser.in
ISSN (Online): 2347-3878
Volume 2 Issue 8, August 2014

Design and Simulation of Optimal Secure Routing Protocol for Wireless Sensor Networks

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Abstract: ***Wireless sensor networks (WSN) represent a new paradigm shift in ad-hoc networks. Wireless Sensor Networks consists of small nodes with sensing, computation and communication capabilities. The sensor networks task is to sense the environment, and relay the information back to a remote base station using multi-hop infrastructure less architecture, where the user can access it. Also since power awareness is an essential design issue in designing a routing protocol for Wireless Sensor Networks, a novel hierarchical based clustering protocol called Efficient Routing Protocol(ERP) for data gathering and aggregation is also proposed for long-lived sensor networks. This proposed protocol handles the redundancy of data problem by using a proper routing tree construction method, thus, ensuring that only effective data is being transmitted through the network. Also it ensures good performance in terms of lifetime by minimizing power consumption for in-network communications and balancing the power load among all the nodes.***

Keywords: sensor, power, energy, base station

1. Introduction

Sensing devices have the capabilities of sensing, computation, self organizing and communication known as sensors. Sensor is a tiny device used to sense the ambient condition of its surroundings, gather data, and process it to draw some meaningful information which can be used to recognize the phenomena around its environment. These sensors can be grouped together using mesh networking protocols to form a network communicating wirelessly using radio frequency channel. The collection of these homogenous or heterogeneous sensor nodes called wireless sensor network (WSN).

2. Sensor Network Architecture

The basic entities in sensor nodes are sensing unit, power unit, processing unit and communication unit and memory unit to perform these operations shown in Figure 2.1 below. Sensor is a hardware device used to measure the change in physical condition of an area of interest and produce response to that change. Sensors sense the environment, collect data and convert it to fundamental data (current or voltage etc) before sending it for further processing. It converts the analogue data (sensed data from an environment) to digital data and then sends it to the microcontroller for further processing. This unit of sensor node is used to store both the data and program code. A node consist a power unit responsible to deliver power to all its units [1]. The basic power consumption at node is due to computation and transmission where transmission is the most expensive activity at sensor node in terms of power consumption. The Processing unit is responsible for data acquisition, processing incoming and outgoing information, implementing and adjusting routing considering the performance conditions of the transmission.

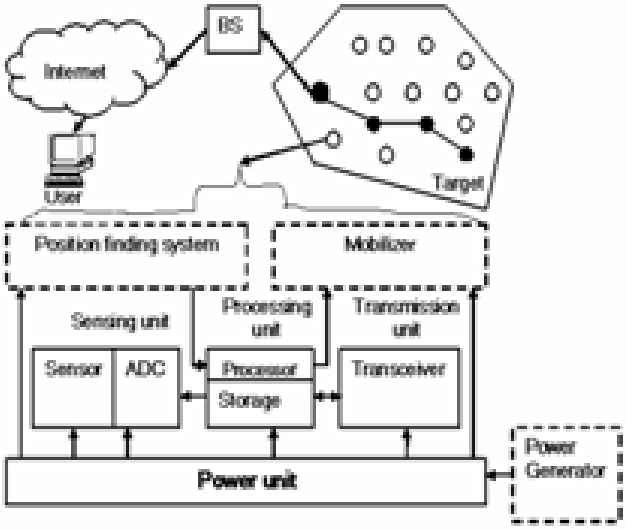

Figure 2.1: Architecture of sensor network

3. Literature Survey and Related Works

In hierarchical-based routing, however, nodes will play different roles in the network. In location-based routing, sensor nodes positions are exploited to route data in the network. A routing protocol is considered adaptive if certain system parameters can be controlled in order to adapt to the current network conditions and available energy levels. The classification is shown in Figure 3.1 where numbers in the Figure indicate the references.

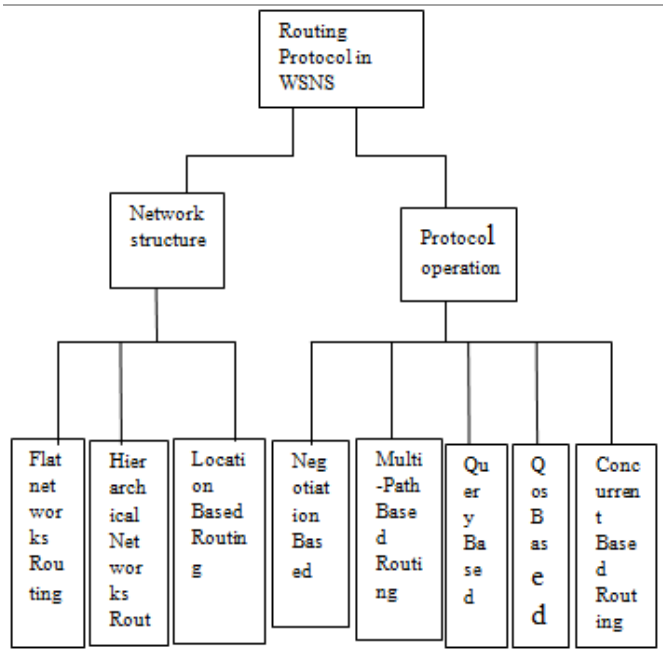

Figure 3.1: Taxonomy of routing protocols in WSN

Hierarchical routing is mainly two-layer routing where one layer is used to select cluster heads and the other layer is used for routing [4].

3.1 LEACH protocol

A hierarchical clustering algorithm for sensor networks, called Low Power Adaptive Clustering Hierarchy (LEACH). LEACH is a cluster-based protocol, which includes distributed cluster formation. LEACH randomly selects a few sensor nodes as cluster heads (CHs) and rotates this role to evenly distribute the power load among the sensors in the network. In LEACH [2], the cluster head (CH) nodes compress data arriving from nodes that belong to the respective cluster, and send an aggregated packet to the base station in order to reduce the amount of information that must be transmitted to the base station.

The operation of LEACH is separated into two phases, the setup phase and the steady state phase. In the setup phase, the clusters are organized and CHs are selected. In the steady state phase, the actual data transfer to the base station takes place. The duration of the steady state phase is longer than the duration of the setup phase in order to minimize overhead. During the setup phase, a predetermined fraction of nodes, p, elect themselves as CHs as follows. A sensor node chooses a random number, r, between 0 and 1. If this random number is less than a threshold value, T (n), the node becomes a cluster-head for the current round.

The threshold value is calculated based on an equation that incorporates the desired percentage to become a cluster-head, the current round, and the set of nodes that have not been selected as a cluster-head in the last (1/P) rounds, denoted by G. It is given by

$$T(n) = \frac{p}{1 - p(r \bmod (1/p))} \quad if\ n \in G$$

where G is the set of nodes that are involved in the CH election. Each elected CH broadcast an advertisement message to the rest of the nodes in the network that they are the new cluster-heads. All the non-cluster head nodes, after receiving this advertisement, decide on the cluster to which they want to belong to. This decision is based on the signal strength of the advertisement. The non cluster-head nodes inform the appropriate cluster-heads that they will be a member of the cluster.

After receiving all the messages from the nodes that would like to be included in the cluster and based on the number of nodes in the cluster, the cluster-head node creates a TDMA schedule and assigns each node a time slot when it can transmit. This schedule is broadcast to all the nodes in the cluster. During the steady state phase, the sensor nodes can begin sensing and transmitting data to the cluster-heads. The cluster-head node, after receiving all the data, aggregates it before sending it to the base-station.

Table 1: Routing protocol comparisons

	Classification	Mobility	Position Awareness	Power Usage	Negotiation based	Data Aggregation	Localization	QoS	State Complexity	Scalability	Multipath	Query based
SPIN	Flat	Possible	No	Limited	Yes	Yes	No	No	Low	Limited	Yes	Yes
Directed Diffusion	Flat	Limited	No	Limited	Yes	Yes	Yes	No	Low	Limited	Yes	Yes
Rumor Routing	Flat	Very Limited	No	N/A	No	Yes	No	No	Low	Good	No	Yes
GBR	Flat	Limited	No	N/A	No	Yes	No	No	Low	Limited	No	Yes
MCFA	Flat	No	No	N/A	No	No	No	No	Low	Good	No	No
CADR	Flat	No	No	Limited	No	Yes	No	No	Low	Limited	No	No
COUGAR	Flat	No	No	Limited	No	Yes	No	No	Low	Limited	No	Yes
ACQUIRE	Flat	Limited	No	N/A	No	Yes	No	No	Low	Limited	No	Yes
EAR	Flat	Limited	No	N/A	No	No		No	Low	Limited	No	Yes
LEACH	Hierarchical	Fixed BS	No	Maximum	No	Yes	Yes	No	CHs	Good	No	No
TEEN & APTEEN	Hierarchical	Fixed BS	No	Maximum	No	Yes	Yes	No	CHs	Good	No	No
PEGASIS	Hierarchical	Fixed BS	No	Maximum	No	No	Yes	No	Low	Good	No	No
MECN & SMECN	Hierarchical	No	No	Maximum	No	No	No	No	Low	Low	No	No
SOP	Hierarchical	No	No	N/A	No	No	No	No	Low	Low	No	No
HPAR	Hierarchical	No	No	N/A	No	No	No	No	Low	Good	No	No
VGA	Hierarchical	No	No	N/A	Yes	Yes	Yes	No	CHs	Good	Yes	No
Sensor aggregate	Hierarchical	Limited	No	N/A	No	Yes	No	No	Low	Good	No	Possible
TTDD	Hierarchical	Yes	Yes	Limited	No	No	No	No	Moderate	Low	Possible	Possible
GAF	Location	Limited	No	Limited	No	No	No	No	Low	Good	No	No
GEAR	Location	Limited	No	Limited	No	No	No	No	Low	Limited	No	No
SPAN	Location	Limited	No	N/A	Yes	No	No	No	Low	Limited	No	No
MFR, GEDIR	Location	No	No	N/A	No	No	No	No	Low	Limited	No	No
GOAFR	Location	No	No	N/A	No	No	No		Low	Good	No	No
SAR	QoS	No	No	N/A	Yes	Yes	No	Yes	Moderate	Limited	No	Yes
SPEED	QoS	No	No	N/A	No	No	No	Yes	moderate	Limited	No	Yes

4. Proposed Approach

Proposed approach presents a general distributed hierarchical clustering scheme called ERP (Efficient Routing Protocol) for saving power consumed by communication. In LEACH, the nodes organize themselves into local clusters, with one node acting as the cluster head. All non-cluster head nodes transmit their data to the cluster head, while the cluster head node receives

the data from all the cluster members, performs signal processing functions on the data (e.g., data aggregation), and transmits data to the remote BS (Base Station). Therefore, being a cluster head is more power intensive than being a non-cluster head node. If the cluster heads were chosen a priori, and fixed throughout the system lifetime, these nodes would quickly use up their limited power. Once the cluster head runs out of power, it is no longer operational. In other words, it is dead. On the other hand, in ERP, after the cluster formation phase, the node with the higher ratio of residual power to average residual power of its cluster range will have a higher probability to become the cluster head.

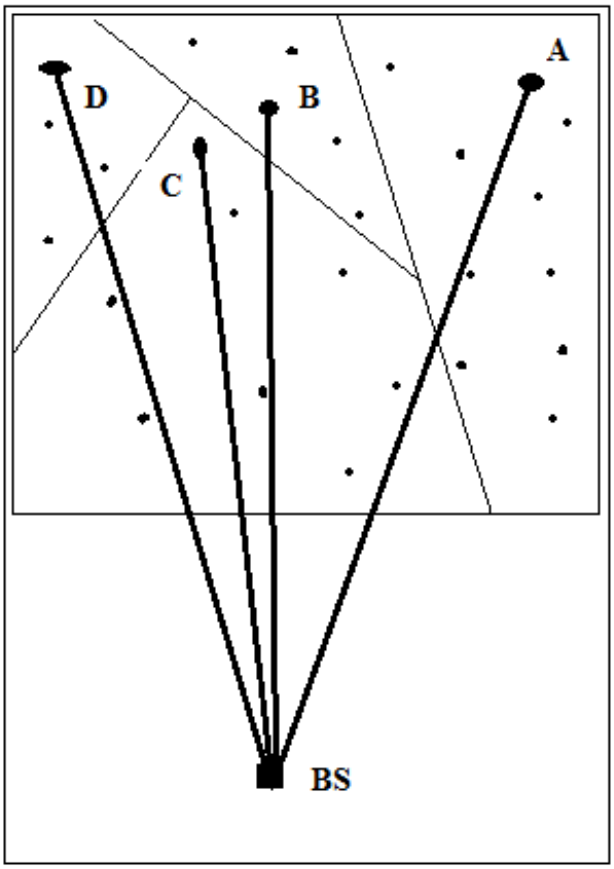

Figure 4.1: Illustration of LEACH Protocol

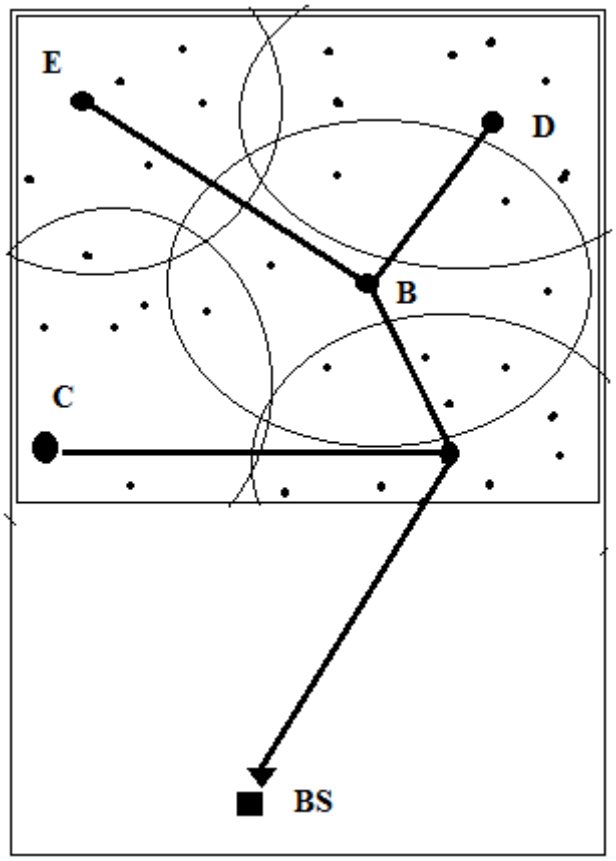

Figure 4.2: Illustration of ERP Protocol

Once, the cluster heads are elected in all the clusters, one node is selected among the cluster heads as the only node (called the root node) to communicate with the BS (also known as the sink node). Thus, all the other cluster heads collect the data from their respective clusters and then send them to the root node. After data aggregation, this data is then sent to the BS as the effective data.

4.1 ERP Protocol

ERP also follows the round operation. Each round includes three phases: $T_{cluster}$, T_{tree}, and T_{data}. In the phase of $T_{cluster}$, after the cluster head selection, the selected cluster head receives the messages sent by all the member nodes that are about to join the cluster head. According to the number of its member nodes, the cluster head creates the TDMA schedule and broadcasts it to all the member nodes in the cluster. In the second phase, T_{tree}, which begins after clustering, one cluster head is selected from the set of other cluster heads as the only node CH_b (Cluster Head at the base or BS) to communicate with the Sink node, so as to increase power efficiency. With CH_b treated as the root, other cluster heads as the offspring of CH_b, are organized to form the appropriate minimum spanning tree.

In the phase of T_{data}, the member node in one cluster, according to the assigned slotted time, sends data to the cluster head. The cluster head continues to wait till it receives the data sent by all its offspring nodes in the spanning tree, and then it sends the aggregated data to its parent node. The above steps are repeated till the aggregated data reaches the BS, which sends the final aggregated data to the sink node.

The time line showing the ERP operation is shown below:

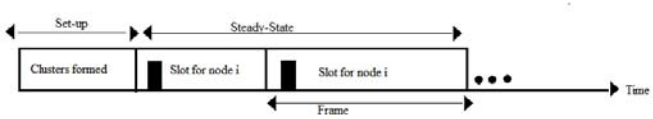

Figure 4.3: Time line showing ERP operation

$T_{cluster}$ and T_{tree} belong to the Set-up phase. T_{data} belongs to the steady-state phase. In the steady-state phase, to reduce power dissipation, each non-cluster head node uses power control to set the amount of transmission power based on the received strength of the cluster head advertisement. Furthermore, the radio of each non-cluster head node is turned off until its allocated transmission time. Since, the design for this approach is going to be optimized when all the nodes have data to send to the cluster head, using a TDMA schedule is an efficient use of bandwidth and represents a low-latency and power-efficient approach.

Pseudo Code of ERP

If(round=1)
Bs_str=Receive_Str("hello")
Broadcast(its E_Msg)
Receive(E_Msg from its neighbour)
Update(neighbourhood table)
If($E_{residual}$<=E_{min})
P=0
Else p=p_{init}
IS_CH=false
While(p<1)
Receive(other messages from its neighbours)
Update(CH_Set)
If(CH_Set=empty&& p<1)
If(random(0,1)<=p)
Broadcast(ID,temp,PRI)
Elseif(CH_Set=empty and p=1)
Broadcast(ID,CH,PRI)
Is_CH=true
Elseif(CH_Set!=empty)
HEAD_ID=Max_PRI(CH_Set)
If(S_i Ɛ CH_Set && S_i'PRI is largest)
If(p=1)
Broadcast(ID,CH,PRI)
IS_CH=true
Else Broadcast(ID,Temp,PRI)
p=min(2*p,1)
If(IS_CH+true)
Broadast(ID,CH)
Else
if(CH_Set=empty)
Broadcast(ID,CH)
else
Join_Cluster(HEAD_ID,ID)

5. Simulation Using OMNet++

5.1 LEACH Protocol

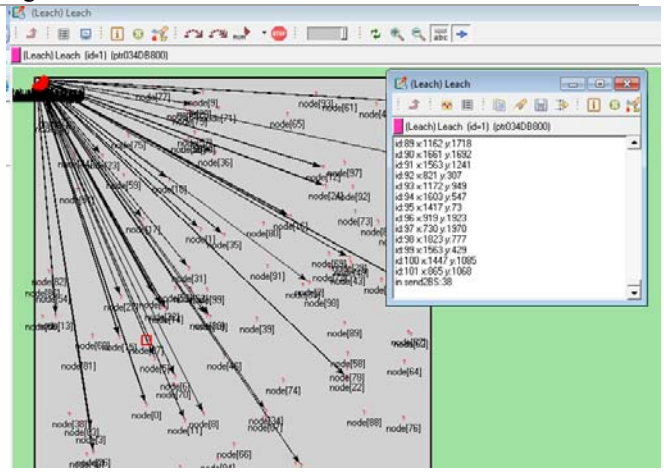

Figure 5.1: Sensor nodes registering with Base station

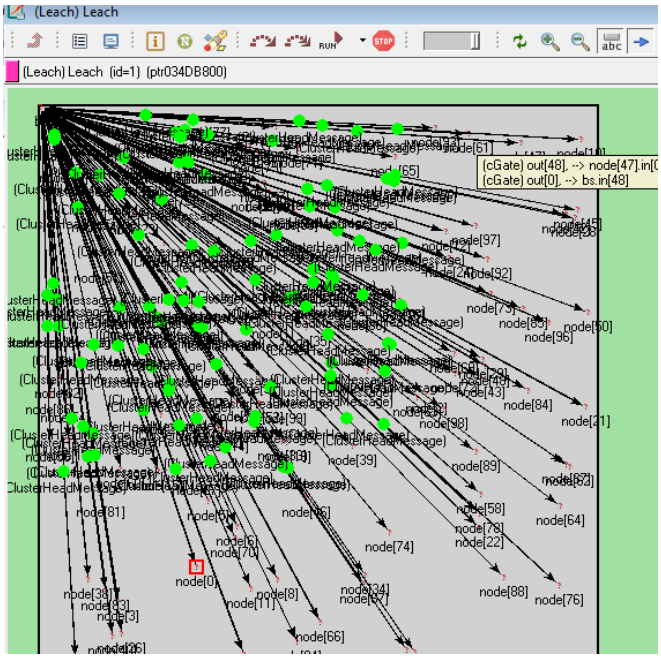

Figure 5.2: Base station send the ACK to Sensor nodes

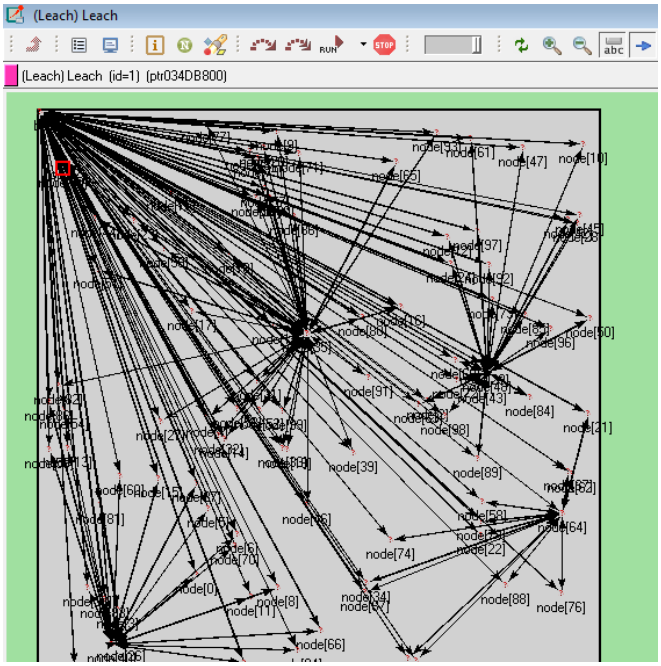

Figure 5.3: LEACH Cluster Formation

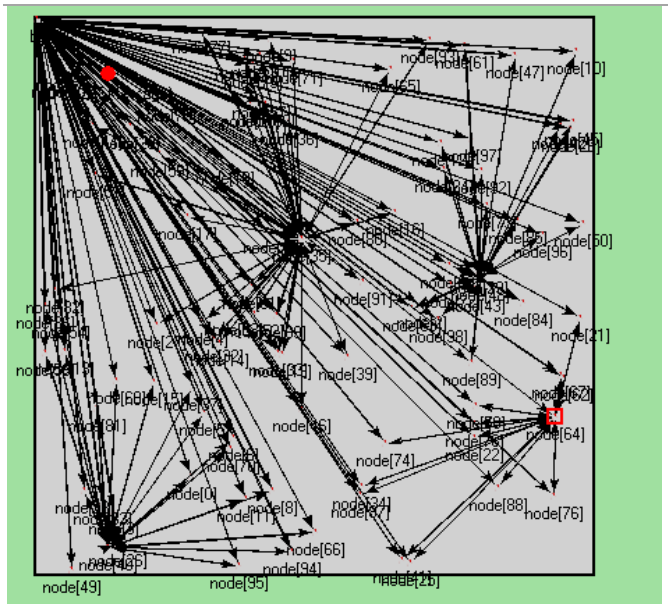

Figure 5.4: Data Transmission phase

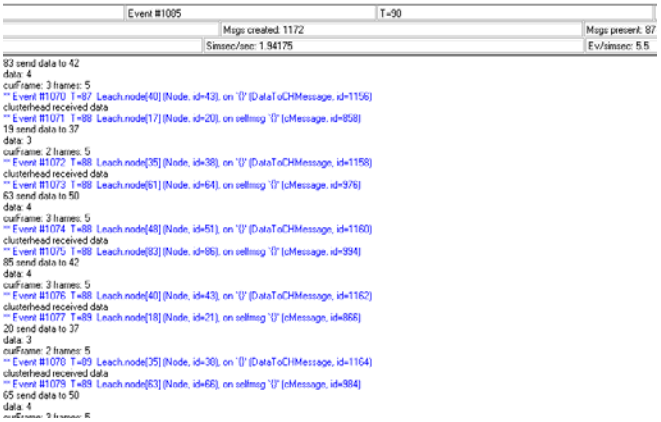

Figure 5.5: Trace file Results for LEACH

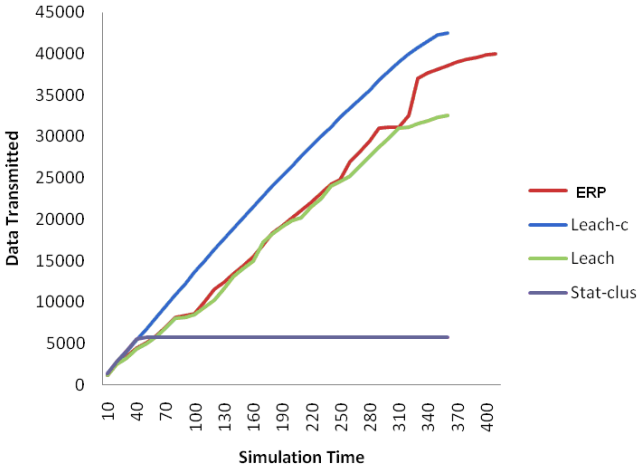

Figure 5.6: Total data received with Time

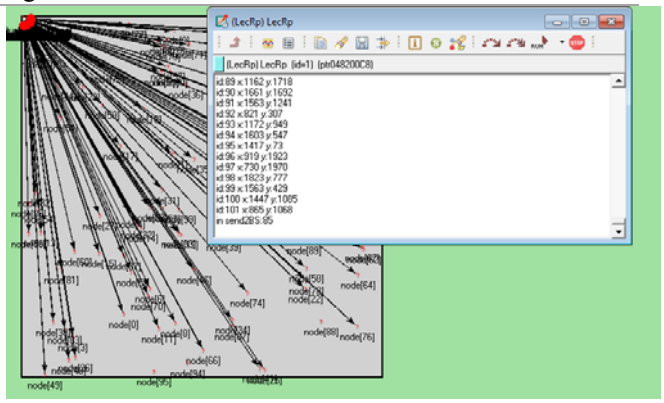

Figure 5.7: Sensor node registering with base station

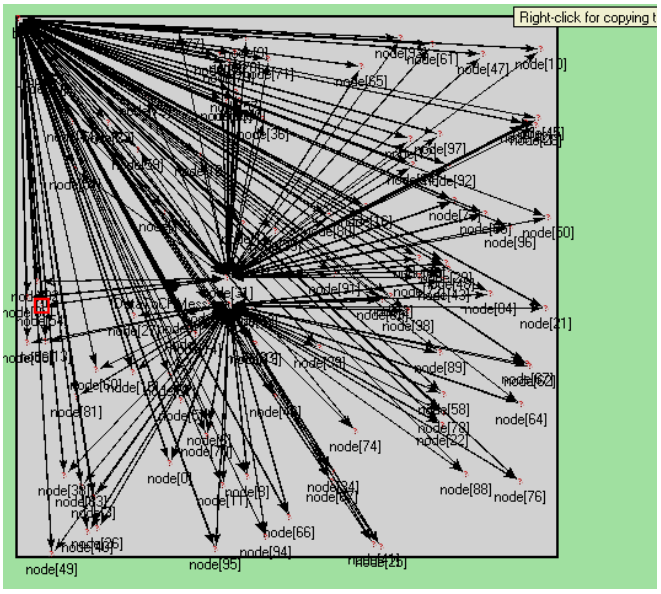

Figure 5.8: Data transmission phase

Figure 5.9: Trace file results for ERP protocol

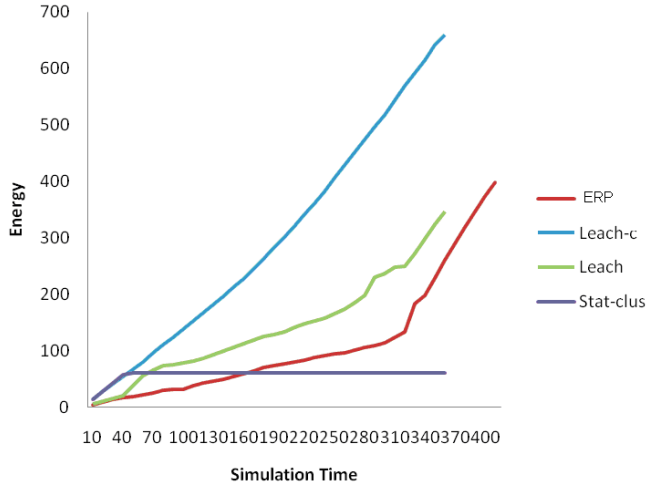

Figure 5.10: Power of the node with Time

From Figure 5.10 which is the illustration of Power, No of nodes alive with respect to simulation-time of the network, we can observe that the power consumed is more in LEACH-C than all the other protocols. ERP performs better than other protocols in terms of power Since ERP protocol lets the power consumption of the network be uniformly distributed in each node, the interval between the deaths of the first node and that of the last node is so short. ERP degrades more slowly than other three protocols in most cases and afterwards it will decrease sharply.

6. Conclusion

In the proposed Efficient Routing Protocol (ERP) for sensor networks, a node independently makes its decision whether or not to compete as a cluster head. After the selection of the cluster heads, the non-cluster head nodes choose to join the closest cluster head and for a cluster. In order to reduce the power consumption, a cluster head, called the base station, is selected from among the other cluster heads as the only node to communicate with the sink node. The cluster head remote from the sink node sends the gathered data, in a multi-hop pattern, to the sink node. Moreover, in ERP protocol, the execution of a simple coverage algorithm helps to prolong the network lifetime with the increase in node density.

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